

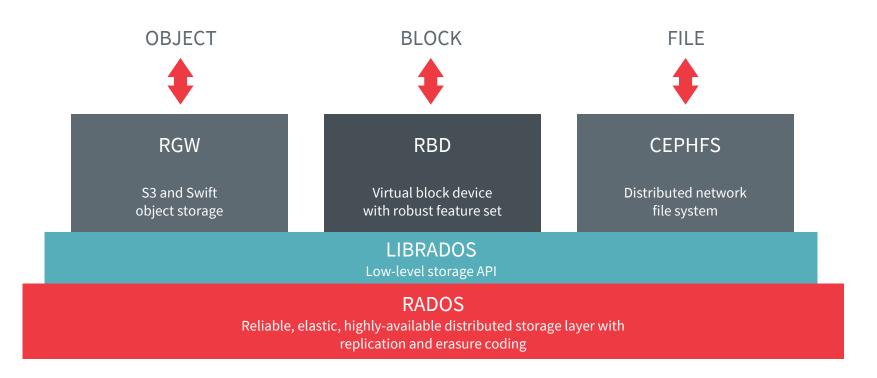
OUTLINE



- Ceph
- Data services
- Block
- File
- Object
- Edge
- Future

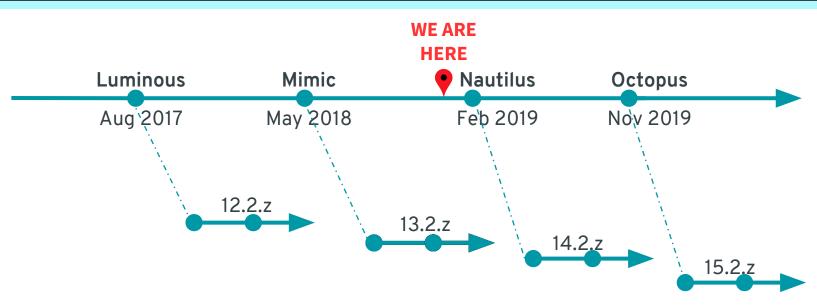
UNIFIED STORAGE PLATFORM





RELEASE SCHEDULE





- Stable, named release every 9 months
- Backports for 2 releases
- Upgrade up to 2 releases at a time
 - (e.g., Luminous → Nautilus, Mimic → Octopus)

FOUR CEPH PRIORITIES



Usability and management

Container platforms

Performance

Multi- and hybrid cloud



A CLOUDY FUTURE



- IT organizations today
 - Multiple private data centers
 - Multiple public cloud services
- It's getting cloudier
 - "On premise" → private cloud
 - Self-service IT resources, provisioned on demand by developers and business units
- Next generation of cloud-native applications will span clouds
- "Stateless microservices" are great, but real applications have state.

DATA SERVICES



Data placement and portability

- Where should I store this data?
- How can I move this data set to a new tier or new site?
- Seamlessly, without interrupting applications?

Introspection

- What data am I storing? For whom? Where? For how long?
- Search, metrics, insights

Policy-driven data management

- Lifecycle management
- Conformance: constrain placement, retention, etc. (e.g., HIPAA, GPDR)
- Optimize placement based on cost or performance
- Automation

MORE THAN JUST DATA



- Data sets are tied to applications
 - When the data moves, the application often should (or must) move too
- Container platforms are key
 - Automated application (re)provisioning
 - "Operators" to manage coordinated migration of state and applications that consume it



DATA USE SCENARIOS



Multi-tier

Different storage for different data

Mobility

- Move an application and its data between sites with minimal (or no) availability interruption
- Maybe an entire site, but usually a small piece of a site

Disaster recovery

- Tolerate a site-wide failure; reinstantiate data and app in a new site quickly
- Point-in-time consistency with bounded latency (bounded data loss)

Stretch

- Tolerate site outage without compromising data availability
- Synchronous replication (no data loss) or async replication (different consistency model)

Edge

Small (e.g., telco POP) and/or semi-connected sites (e.g., autonomous vehicle)



BLOCK STORAGE



HOW WE USE BLOCK



- Virtual disk device
- Exclusive access by nature (with few exceptions)
- Strong consistency required
- Performance sensitive
- Basic feature set
 - Read, write, flush, maybe resize
 - Snapshots (read-only) or clones (read/write)
 - Point-in-time consistent
- Often self-service provisioning
 - via Cinder in OpenStack
 - o via Persistent Volume (PV) abstraction in Kubernetes

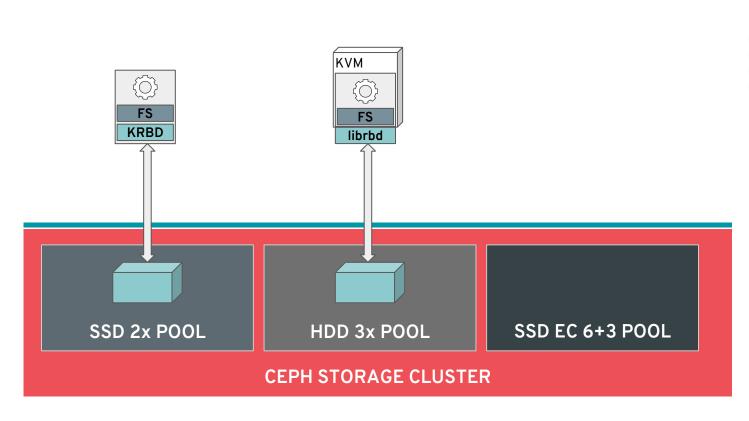
Applications

XFS, ext4, whatever

Block device

RBD - TIERING WITH RADOS POOLS

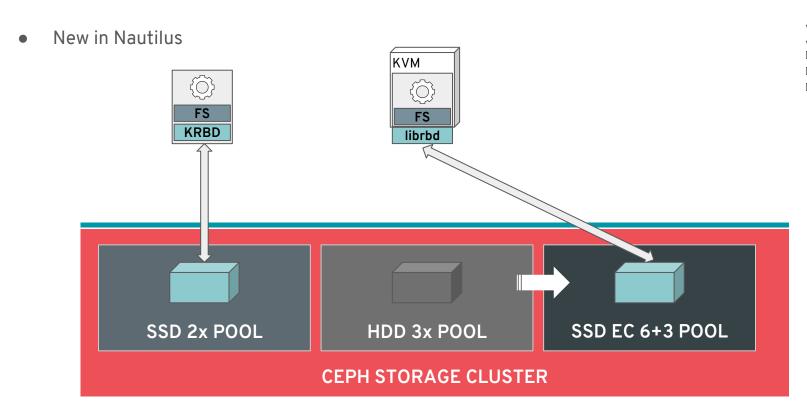




- Mobility DR
- □ Stretch
- **E**dge

RBD - LIVE IMAGE MIGRATION

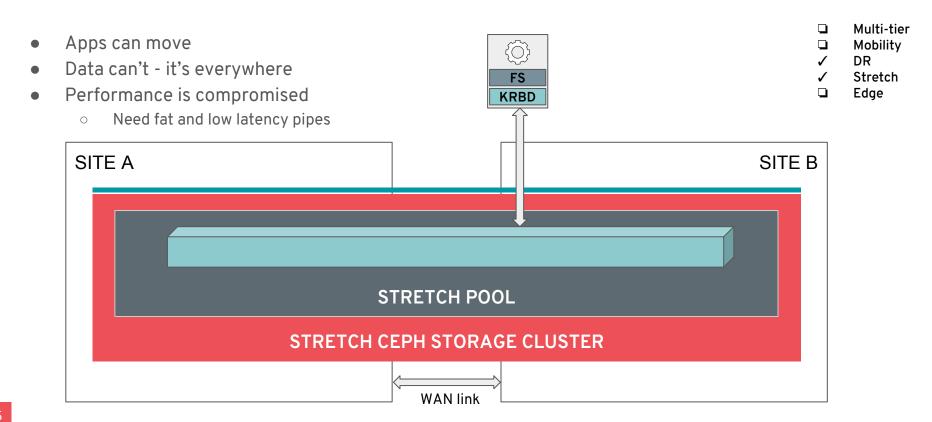




- ✓ Mobility
- **□** DR
- ⊒ Stretch ⊒ Edge

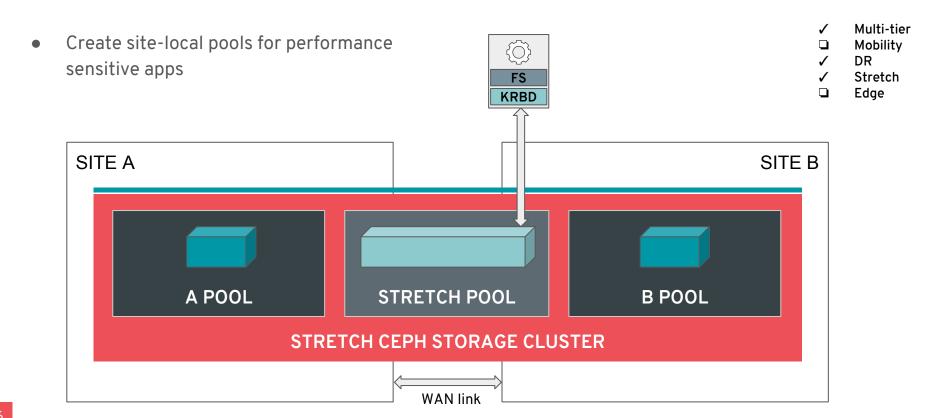
RBD - STRETCH





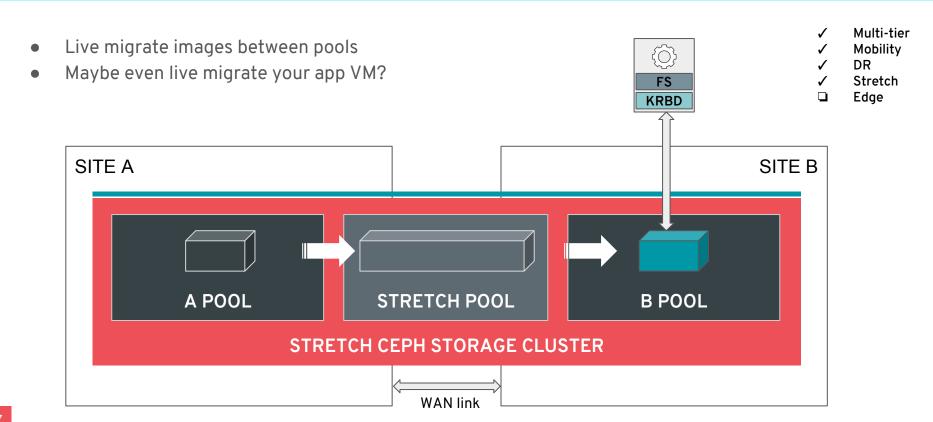
RBD - STRETCH WITH TIERS





RBD - STRETCH WITH MIGRATION





STRETCH IS SKETCH

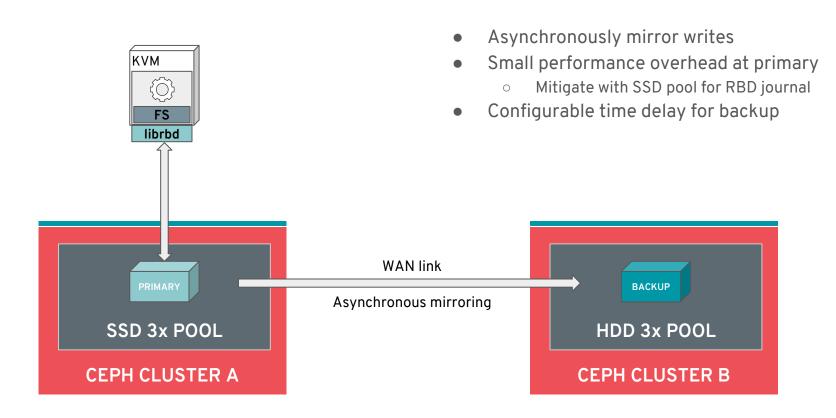


- Network latency is critical
 - Low latency for performance
 - Requires nearby sites, limiting usefulness
- Bandwidth too
 - Must be able to sustain rebuild data rates
- Relatively inflexible
 - Single cluster spans all locations
 - Cannot "join" existing clusters
- High level of coupling
 - Single (software) failure domain for all sites



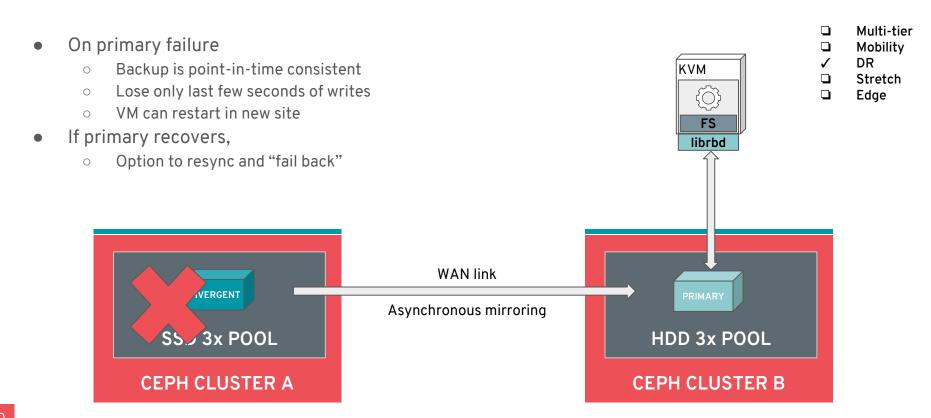
RBD ASYNC MIRRORING





RBD ASYNC MIRRORING





RBD MIRRORING IN CINDER



- Ocata
 - Cinder RBD replication driver
- Queens
 - ceph-ansible deployment of rbd-mirror via TripleO
- Rocky
 - Failover and fail-back operations

Gaps

- Deployment and configuration tooling
- o Cannot replicate multi-attach volumes
- Nova attachments are lost on failover



MISSING LINK: APPLICATION ORCHESTRATION



- Hard for laaS layer to reprovision app in new site
- Storage layer can't solve it on its own either
- Need automated, declarative, structured specification for entire app stack...





FILE STORAGE



CEPHFS STATUS



- Stable since Kraken
- Multi-MDS stable since Luminous
- Snapshots stable since Mimic
- Support for multiple RADOS data pools
- Provisioning via OpenStack Manila and Kubernetes
- Fully awesome

- ✓ Multi-tier □ Mobility
- DR
- ⊒ Stretch
- Edge

CEPHFS - STRETCH?



- We can stretch CephFS just like RBD pools
- It has the same limitations as RBD
 - Latency → lower performance
 - Limited by geography
 - o Big (software) failure domain
- Also,
 - MDS latency is critical for file workloads
 - o ceph-mds daemons be running in one site or another

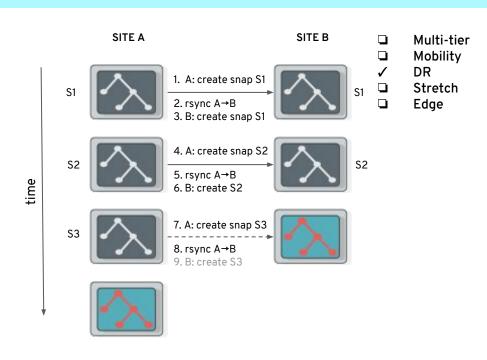
What can we do with CephFS across multiple clusters?

- ☐ Multi-tier ☐ Mobility
- ✓ DR
- ✓ Stretch
 - Edge

CEPHFS - SNAP MIRRORING



- CephFS snapshots provide
 - o point-in-time consistency
 - granularity (any directory in the system)
- CephFS rstats provide
 - o rctime to efficiently find changes
- rsync provides
 - efficient file transfer
- Time bounds on order of minutes
- Gaps and TODO
 - "rstat flush" coming in Nautilus
 - Xuehan Xu @ Qihoo 360
 - rsync support for CephFS rstats
 - scripting / tooling



DO WE NEED POINT-IN-TIME FOR FILE?



- Yes.
- Sometimes.
- Some geo-replication DR features are built on rsync...
 - Consistent view of individual files,
 - Lack point-in-time consistency between files
- Some (many?) applications are not picky about cross-file consistency...
 - Content stores
 - Casual usage without multi-site modification of the same files

CASE IN POINT: HUMANS



- Many humans love Dropbox / NextCloud / etc.
 - Ad hoc replication of directories to any computer
 - Archive of past revisions of every file
 - Offline access to files is extremely convenient and fast
- Disconnected operation and asynchronous replication leads to conflicts
 - Usually a pop-up in GUI
- Automated conflict resolution is usually good enough
 - o e.g., newest timestamp wins
 - Humans are happy if they can rollback to archived revisions when necessary
- A possible future direction:
 - Focus less on avoiding/preventing conflicts...
 - Focus instead on ability to rollback to past revisions...

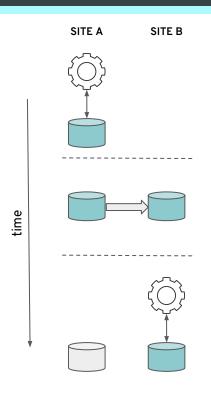
BACK TO APPLICATIONS



- Do we need point-in-time consistency for file systems?
- Where does the consistency requirement come in?

MIGRATION: STOP, MOVE, START





- App runs in site A
- Stop app in site A
- Copy data A→B
- Start app in site B

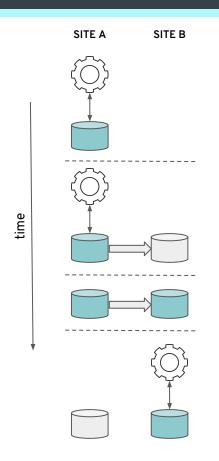
- ☐ Multi-tier
- ✓ Mobility

 □ DR
- □ Stretch □ Edge

- App maintains exclusive access
- Long service disruption

MIGRATION: PRESTAGING



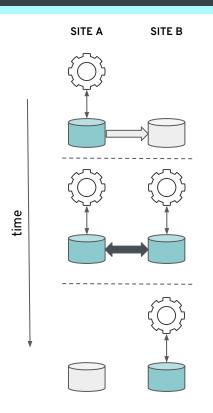


- App runs in site A
- Copy most data from A→B
- Stop app in site A
- Copy last little bit A→B
- Start app in site B

- App maintains exclusive access
- Short availability blip

MIGRATION: TEMPORARY ACTIVE/ACTIVE





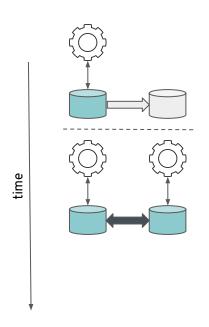
- App runs in site A
- Copy most data from A→B
- Enable bidirectional replication
- Start app in site B
- Stop app in site A
- Disable replication

- No loss of availability
- Concurrent access to same data

ACTIVE/ACTIVE



SITE A SITE B



- App runs in site A
- Copy most data from A→B
- Enable bidirectional replication
- Start app in site B

- Highly available across two sites
- Concurrent access to same data

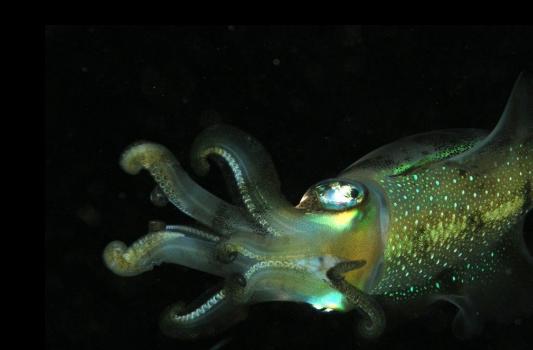
BIDIRECTIONAL FILE REPLICATION?



- We don't have general-purpose bidirectional file replication
- It is hard to resolve conflicts for any POSIX operation
 - Sites A and B both modify the same file
 - Site A renames $/a \rightarrow /b/a$ while Site B: renames $/b \rightarrow /a/b$
- But applications can only go active/active if they are cooperative
 - i.e., they carefully avoid such conflicts
 - e.g., mostly-static directory structure + last writer wins
- So we could do it if we simplify the data model...
- But wait, that sounds a bit like object storage...



OBJECT STORAGE



WHY IS OBJECT SO GREAT?



Based on HTTP

Interoperates well with web caches, proxies, CDNs, ...

Atomic object replacement

- PUT on a large object atomically replaces prior version
- Trivial conflict resolution (last writer wins)
- Lack of overwrites makes erasure coding easy

Flat namespace

- No multi-step traversal to find your data
- Easy to scale horizontally

No rename

Vastly simplified implementation

THE FUTURE IS... OBJECTY



- File is not going away, and will be critical
 - Half a century of legacy applications
 - It's genuinely useful
- Block is not going away, and is also critical infrastructure
 - Well suited for exclusive-access storage users (boot devices, etc)
 - Performs better than file due to local consistency management, ordering etc.
- Most new data will land in objects
 - o Cat pictures, surveillance video, telemetry, medical imaging, genome data
 - Next generation of cloud native applications will be architected around object

RGW FEDERATION MODEL



Zone

- Collection RADOS pools storing data
- Set of RGW daemons serving that content

ZoneGroup

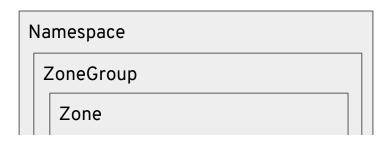
- Collection of Zones with a replication relationship
- Active/Passive[/...] or Active/Active

Namespace

- Independent naming for users and buckets
- All ZoneGroups and Zones replicate user and bucket index pool
- One Zone serves as the leader to handle User and Bucket creations/deletions

Failover is driven externally

 Human (?) operators decide when to write off a master, resynchronize

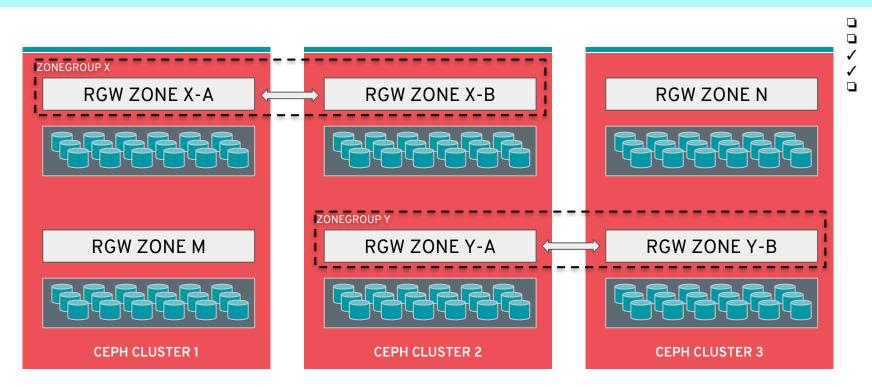


RGW FEDERATION TODAY



Multi-tier Mobility DR

Stretch Edge



• Gap: granular, per-bucket management of replication

ACTIVE/ACTIVE FILE ON OBJECT

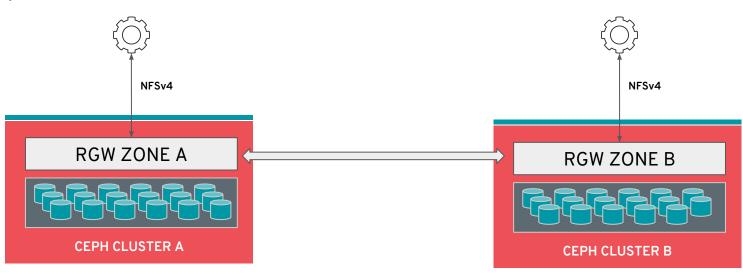


DR Stretch

Edge

Multi-tier Mobility

- Data in replicated object zones
 - Eventually consistent, last writer wins
- Applications access RGW via NFSv4



OTHER RGW REPLICATION PLUGINS



ElasticSearch

- Index entire zone by object or user metadata
- Query API

Cloud sync (Mimic)

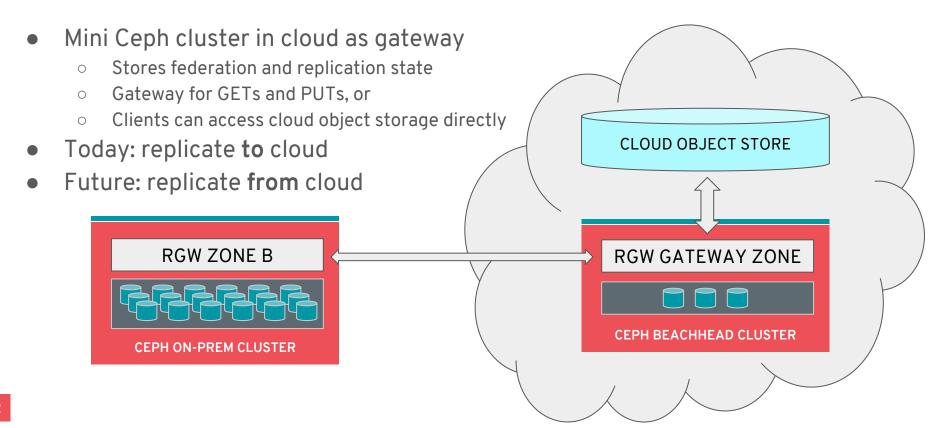
- Replicate buckets to external object store (e.g., S3)
- Can remap RGW buckets into multiple S3 bucket, same S3 bucket
- Remaps ACLs, etc

Archive (Nautilus)

- Replicate all writes in one zone to another zone, preserving all versions
- Pub/sub (Nautilus)
 - Subscribe to event notifications for actions like PUT
 - Integrates with knative serverless! (See Huamin and Yehuda's talk at Kubecon next month)

PUBLIC CLOUD STORAGE IN THE MESH





RGW TIERING



DR Stretch Edge

Multi-tier Mobility

<u>Today: Intra-cluster</u>

- Many RADOS pools for a single RGW zone
- Primary RADOS pool for object "heads"
 - Single (fast) pool to find object metadata and location of the tail of the object
- Each tail can go in a different pool
 - Specify bucket policy with PUT
 - Per-bucket policy as default when not specified
- Policy
 - Retention (auto-expire)

<u>Nautilus</u>

- Tier objects to an external store
 - Initially something like S3
 - Later: tape backup, other backends...

Later

- Encrypt data in external tier
- Compression
- (Maybe) cryptographically shard across multiple backend tiers
- Policy for moving data between tiers

RGW - THE BIG PICTURE



<u>Today</u>

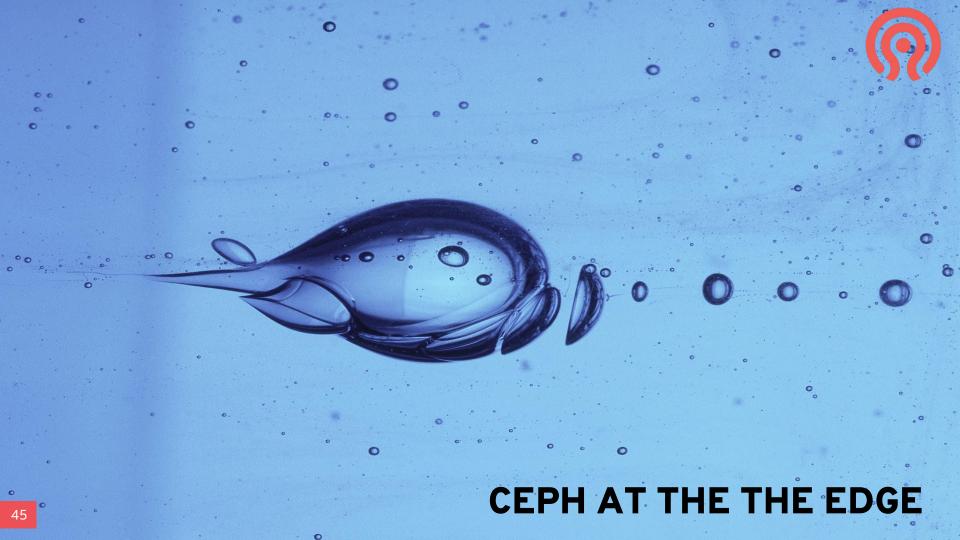
- RGW as gateway to a RADOS cluster
 - With some nifty geo-replication features
- RGW redirects clients to the correct zone
 - via HTTP Location: redirect
 - Dynamic DNS can provide right zone IPs

- RGW replicates at zone granularity
 - Well suited for disaster recovery

Future

- RGW as a gateway to a mesh of sites
 - With great on-site performance
- RGW may redirect or proxy to right zone
 - Single point of access for application
 - Proxying enables coherent local caching

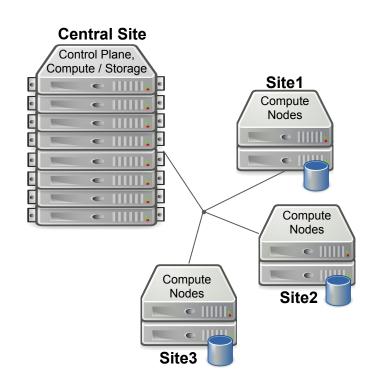
- RGW may replicate at bucket granularity
 - Individual applications set durability needs
 - Enable granular application mobility



CEPH AT THE EDGE



- A few edge examples
 - Telco POPs: ¼ ½ rack of OpenStack
 - Autonomous vehicles: cars or drones
 - Retail
 - Backpack infrastructure
- Scale down cluster size
 - Hyper-converge storage and compute
 - Nautilus: brings better memory control
- Multi-architecture support
 - o aarch64 (ARM) builds upstream
 - o POWER builds at OSU / OSL
- Hands-off operation
 - Ongoing usability work
 - Operator-based provisioning (Rook)
- Possibly unreliable WAN links



DATA AT THE EDGE



- Block: async mirror edge volumes to central site
 - For DR purposes
- Data producers
 - Write generated data into objects in local RGW zone
 - Upload to central site when connectivity allows
 - Perhaps with some local pre-processing first
- Data consumers
 - Access to global data set via RGW (as a "mesh gateway")
 - Local caching of a subset of the data
- We're most interested in object-based edge scenarios

_	Multi-tie
ב	Mobility
ב	DR
]	Stretch
/	Eda.



WHY ALL THE KUBERNETES TALK?







- True mobility is a partnership between orchestrator and storage
- Kubernetes is an emerging leader in application orchestration
- Persistent Volumes
 - Basic Ceph drivers in Kubernetes, ceph-csi on the way
 - o Rook for automating Ceph cluster deployment and operation, hyperconverged

Object

- Trivial provisioning of RGW via Rook
- Coming soon: on-demand, dynamic provisioning of Object Buckets and User (via Rook)
- Consistent developer experience across different object backends (RGW, S3, minio, etc.)



BRINGING IT ALL TOGETHER...

SUMMARY



- Data services: mobility, introspection, policy
- These are a partnership between storage layer and application orchestrator
- Ceph already has several key multi-cluster capabilities
 - Block mirroring
 - Object federation, replication, cloud sync; cloud tiering, archiving and pub/sub coming
 - Cover elements of Tiering, Disaster Recovery, Mobility, Stretch, Edge scenarios
- ...and introspection (elasticsearch) and policy for object
- Future investment is primarily focused on object
 - RGW as a gateway to a federated network of storage sites
 - Policy driving placement, migration, etc.
- Kubernetes will play an important role
 - both for infrastructure operators and applications developers

