

UNLOCK BIGDATA ANALYTIC EFFICIENCY WITH CEPH DATA LAKE

Kyle Bader - Red Hat

Yuan Zhou, Yong Fu, Jian Zhang - Intel

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Agenda

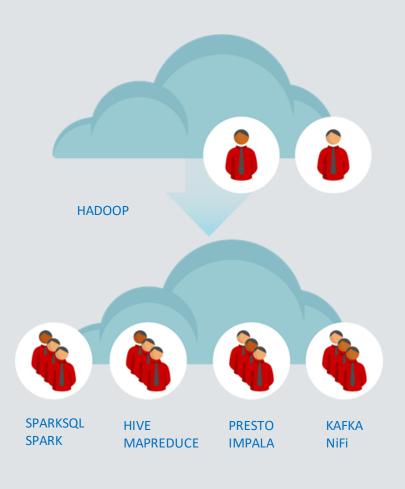
- Background and Motivations
- The Workloads, Reference Architecture Evolution and Performance Optimization
- Performance Comparison with Remote HDFS
- Summary & Next Step



BACKGROUND AND MOTIVATION



DISCONTINUITY IN BIG DATA INFRASTRUCTURE - WHY?



CONGESTION

in busy analytic clusters causing missed SLAs.

MULTIPLE TEAMS COMPETING

and sharing the same big data resources.

MODERN BIG DATA ANALYTICS PIPELINE KEY TERMINOLOGY















MODERN BIG DATA ANALYTICS PIPELINE KEY TERMINOLOGY



Kafka



- Hadoop
- Spark



- Spark
- Hadoop



- Sensors
- Click-stream
- Transactions
- Call-detail records



- NiFi
- Kafka

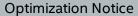


DATA SCIENCE

- Presto
- Impala
- SparkSQL

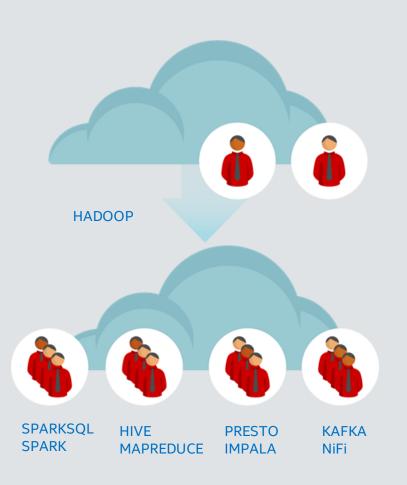


TensorFlow





DISCONTINUITY IN BIG DATA INFRASTRUCTURE - WHY?



CONGESTION

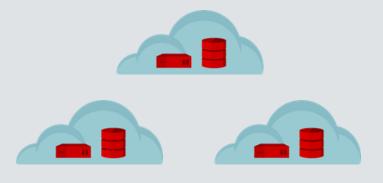
in busy analytic clusters causing missed SLAs.

MULTIPLE TEAMS COMPETING

and sharing the same big data resources.

CAUSING CUSTOMERS TO PICK A SOLUTION







#1

Get a bigger cluster for many teams to share.

#2

Give each team their own dedicated cluster, each with a copy of PBs of data.

#3

Give teams ability to spin-up/spin-down clusters which can share data sets.





#1 SINGLE LARGE CLUSTER



- Lacks isolation noisy neighbors hinder SLAs.
- Lacks elasticity single rigid cluster.



#2 MULTIPLE SMALL CLUSTERS



- No dataset sharing
- Cost of duplicate storage
- Still lacks elasticity
- Can't scale



#3 ON DEMAND ANALYTIC CLUSTERS WITH A SHARED DATA LAKE





HIT SERVICE-LEVEL AGREEMENTS

Give teams their own compute clusters.



ELIMINATE IDLE RESOURCES

By right-sizing de-coupled compute and storage.



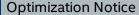
BUY 10s OF PBS INSTEAD OF 100S

Share data sets across clusters instead of duplicating them.



INCREASE AGILITY

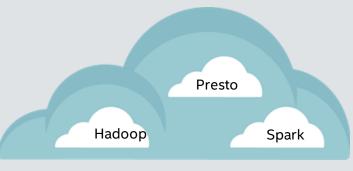
With spin-up/spin-down clusters.





#3 ON DEMAND ANALYTIC CLUSTERS WITH A SHARED DATA LAKE

PUBLIC CLOUD (AWS)



AWS EC2
PROVISIONING



AWS S3
SHARED DATASETS

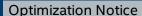
PRIVATE CLOUD



OPENSTACK PROVISIONING



CEPH SHARED DATASETS





GENERATION 1MONOLITHIC HADOOP STACKS

CIOUDETA

HORTONWORKS*
POWERING THE FUTURE OF DATA*

ANALYTICS +
INFRASTRUCTURE

Analytics vendors provide single-purpose infrastructure

___ Analytics vendors provide analytics software



GENERATION 2

DECOUPLED STACK WITH PRIVATE CLOUD INFRASTRUCTURE











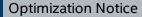


Analytics vendors provide analytics software.

Provisioned Compute Pool via OpenStack

Shared Datasets on Ceph Object Store

Private cloud provides
Infrastructure services





THE WORKLOADS, REFERENCE ARCHITECTURE AND PERFORMANCE



Workloads

Simple Read/Write

- DFSIO: TestDFSIO is the canonical example of a benchmark that attempts to measure the Storage's capacity for reading and writing bulk data.
- Terasort: a popular benchmark that measures the amount of time to sort one terabyte of randomly distributed data on a given computer system.

Data Transformation

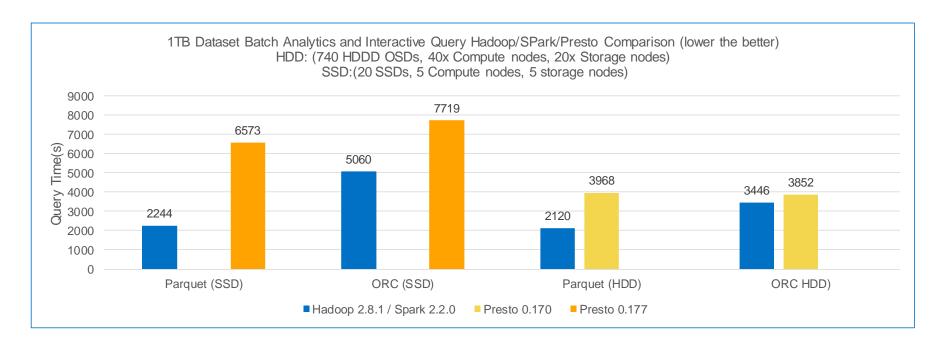
 ETL: Taking data as it is originally generated and transforming it to a format (Parquet, ORC) that more tuned for analytical workloads.

Batch Analytics

- To consistently executing analytical process to process large set of data.
- Leveraging 54 derived from TPC-DS * queries with intensive reads across objects in different buckets



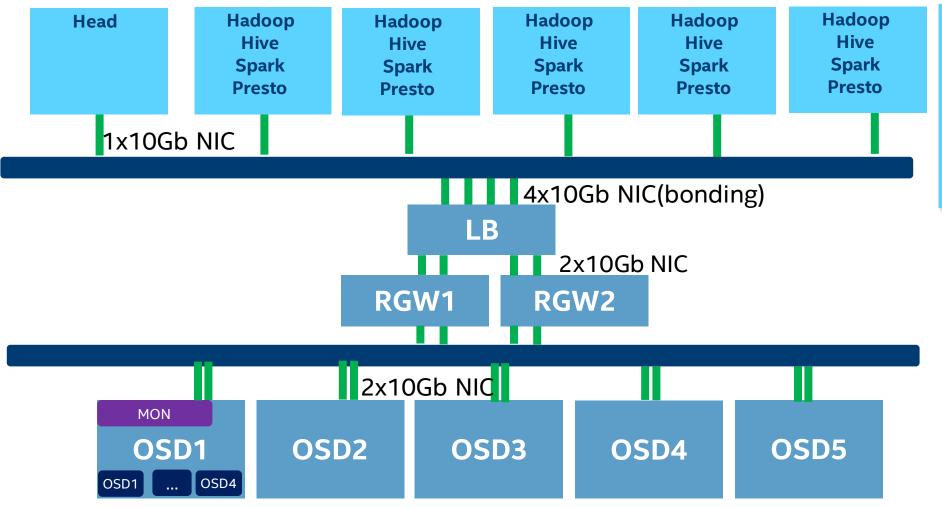
Bigdata on Object Storage Performance Overview --Batch analytics



- Significant performance improvement from Hadoop 2.7.3/Spark 2.1.1 to Hadoop 2.8.1/Spark 2.2.0 (improvement in s3a)
- Batch analytics performance of 10-node Intel AFA is almost on-par with 60-node HDD cluster



Hardware Configuration --Dedicate LB



5x Compute Node

- Intel® Xeon™ processor E5-2699 v4 @ 2.2GHz, 128GB mem
- 2x10G 82599 10Gb NIC
- 2x SSDs
- 3x Data storage (can be emliminated)

Software:

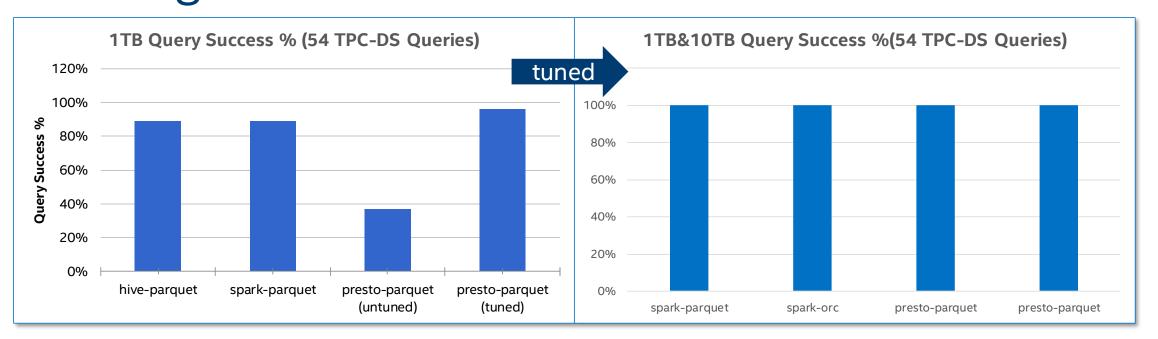
- Hadoop 2.7.3
- Spark 2.1.1
- Hive 2.2.1
- Presto 0.177
- RHEL7.3

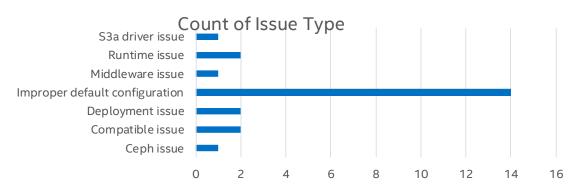
5x Storage Node, 2 RGW nodes, 1 LB nodes

- Intel(R) Xeon(R) CPU E5-2699v4 2.20GHz
- 128GB Memory
- 2x 82599 10Gb NIC
- 1x Intel® P3700 1.0TB SSD as Journal
- 4x 1.6TB Intel® SSD DC S3510 as data drive
- 2x 400G S3700 SSDs
- 1 OSD instances one each S3510 SSD
- RHEl7.3
- RHCS 2.3



Improve Query Success Ratio with Functional Trouble-shooting





- 100% selected TPC-DS query passed with tunings
- Improper Default configuration
 - small capacity size,
 - wrong middleware configuration
 - improper Hadoop/Spark configuration for different size and format data issues



Optimizing HTTP Requests -- The bottlenecks

2017-07-18 14:53:52.2599767fddd67fc700 1 ====== starting new request req=0x7fddd67f6710 ===== 2017-07-18 14:53:52.2718297fddd5ffb700 1 ====== starting new request req=0x7fddd5ff5710 ===== 2017-07-18 14:53:52.2739407fddd7fff700 0 ERROR: flush_read_list(): d->client_c->handle_data() returned - 5

2017-07-18 14:53:52.2742237fddd7fff700 0 WARNING: set_req_state_err err_no=5 resorting to 500 2017-07-18 14:53:52.2742537fddd7fff700 0 ERROR: s->cio->send_content_length() returned err=-5 2017-07-18 14:53:52.2742577fddd7fff700 0 ERROR: s->cio->print() returned err=-5 2017-07-18 14:53:52.2742587fddd7fff700 0 ERROR: STREAM_IO(s)->print() returned err=-5 2017-07-18 14:53:52.2742677fddd7fff700 0 ERROR: STREAM_IO(s)->complete header() returned err=-5

Http 500 errors in RGW log **ESTAB**

ESTAB

ESTAB

0

0 0

0 0

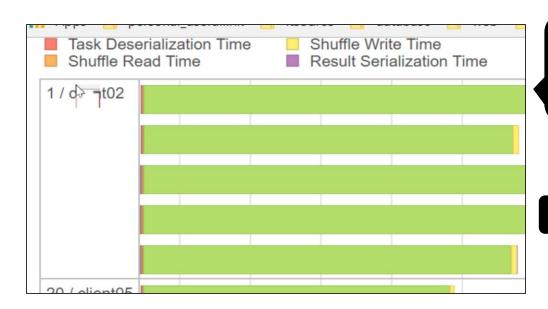
0

0 0

0

ESTAB

ESTAB



Compute time take the big part. (compute time = read data +sort)

New connections out every time,
Connection not reused
ESTAB

```
ESTAB
                        1597240
                                        ::ffff:10.0.2.36:44436
                            ::ffff:10.0.2.254:80
                       ::ffff:10.0.2.3<mark>6:44448</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
ESTAB
                       ::ffff:10.0.2.3<mark>6:44338</mark>
           0
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                       ::ffff:10.0.2.3<mark>6:44438</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                       ::ffff:10.0.2.3<mark>6:44414</mark>
                                                        ::ffff:10.0.2.254:80
                0
                        ::ffff:10.0.2.36:44450
                                                         ::ffff:10.0.2.254:80
ESTAB
          0
               480
                           timer:(on,170ms,0)
                       ::ffff:10.0.2.3<mark>6:44442</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
ESTAB
           0
                       ::ffff:10.0.2.3<mark>6:44390</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                       ::ffff:10.0.2.3<mark>6:44326</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                0
                       ::ffff:10.0.2.3<mark>6:44452</mark>
                                                        ::ffff:10.0.2.254:80
                       ::ffff:10.0.2.3<mark>6:44394</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
ESTAB
           0
                       ::ffff:10.0.2.3<mark>6:44444</mark>
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                0
                       ::ffff:10.0.2.36:44456
                                                        ::ffff:10.0.2.254:80
            2 seconds interval =========
                                                        ::ffff:10.0.2.254:80
ESTAB
                       ::ffff:10.0.2.36:44508
ESTAB
           0
                0
                       ::ffff:10.0.2.36:44476
                                                        ::ffff:10.0.2.254:80
ESTAB
                       ::ffff:10.0.2.36:44524
           0
                                                        ::ffff:10.0.2.254:80
ESTAB
                       ::ffff:10.0.2.36:44374
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                0
                       ::ffff:10.0.2.36:44500
                                                        ::ffff:10.0.2.254:80
                       ::ffff:10.0.2.36:44504
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
ESTAB
           0
                       ::ffff:10.0.2.36:44512
                                                        ::ffff:10.0.2.254:80
ESTAB
           0
                0
                       ::ffff:10.0.2.36:44506
                                                        ::ffff:10.0.2.254:80
           0
                       ::ffff:10.0.2.36:44464
                                                        ::ffff:10.0.2.254:80
                       ::ffff:10.0.2.36:44518
                                                        ::ffff:10.0.2.254:80
```

::ffff:10.0.2.36:44510

::ffff:10.0.2.36:44442

::ffff:10.0.2.36:44526

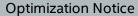
::ffff:10.0.2.36:44472

::ffff:10.0.2.36:44466

::ffff:10.0.2.36.44446

::ffff:10.0.2.3<mark>6:44454</mark>

::ffff:10.0.2.3<mark>6:44374</mark>





::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

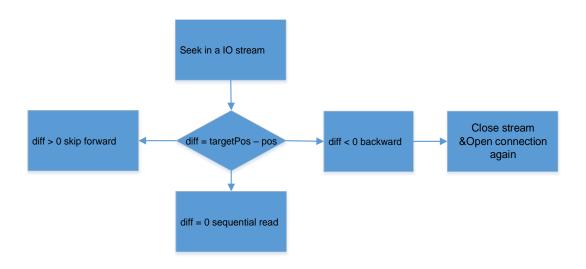
::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

::ffff:10.0.2.254:80

Optimizing HTTP Requests -- S3a input policy



Background

The S3A filesystem client supports the notion of input policies, similar to that of the POSIX fadvise() API call. This tunes the behavior of the S3A client to optimize HTTP GET requests for various use cases. To optimize HTTP GET requests, you can take advantage of the S3A experimental input policy fs.s3a.experimental.input.fadvise.

Ticket: https://issues.apache.org/jira/browse/HADOOP-13203

Solution

Enable random read policy hadoop:

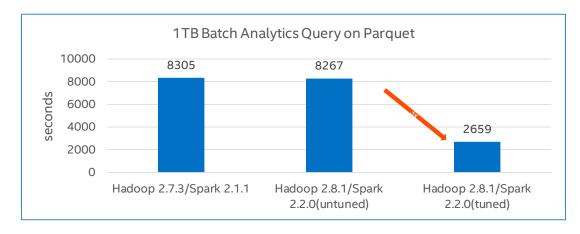
```
<property>
    <name>fs.s3a.experimental.input.fadvise</name>
    <value>random</value>
    </property>
    <property>
    <name>fs.s3a.readahead.range</name>
         <value>64K</value>
         </property></property>
```

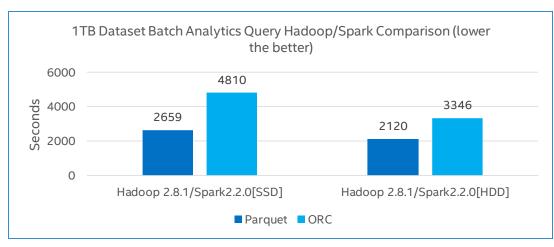
By reducing the cost of closing existing HTTP requests, this is highly efficient for file IO accessing a binary file through a series of `PositionedReadable.read()` and `PositionedReadable.readFully()` calls.



Optimizing HTTP Requests

-- Performance

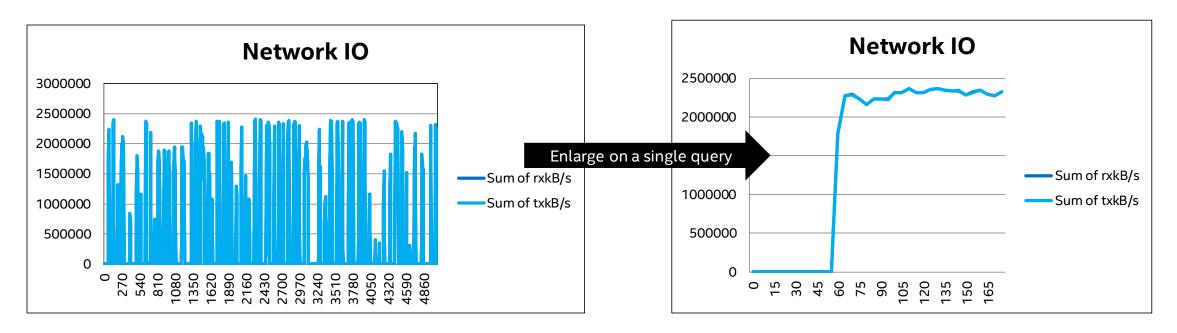




- Readahead feature is supported from Hadoop 2.8.1, but not enabled by default. By applying random read policy, the 500 issue is fixed and performance improved 3x
- All Flash storage architecture also show great performance benefit and low TCO which compared with HDD storage



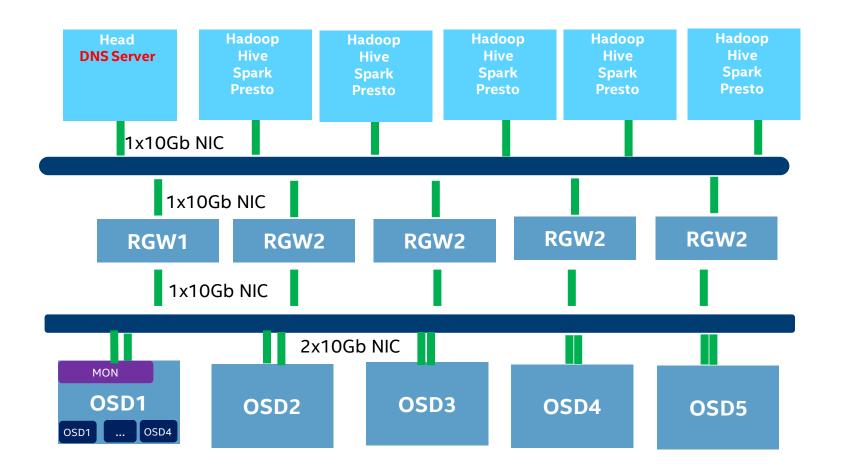
New bottleneck on Load Balancer



- Load Balancer became the bottleneck on networking bandwidth
- Observed many messages blocked at load balancer server(send to s3a driver), but not much blocked at receiving on s3a driver side



Hardware Configuration --More RGWs with round-robin DNS



5x Compute Node

- Intel® Xeon™ processor E5-2699 v4 @ 2.2GHz, 128GB mem
- 2x10G 82599 10Gb NIC
- 2x SSDs
- 3x Data storage (can be emliminated)

Software:

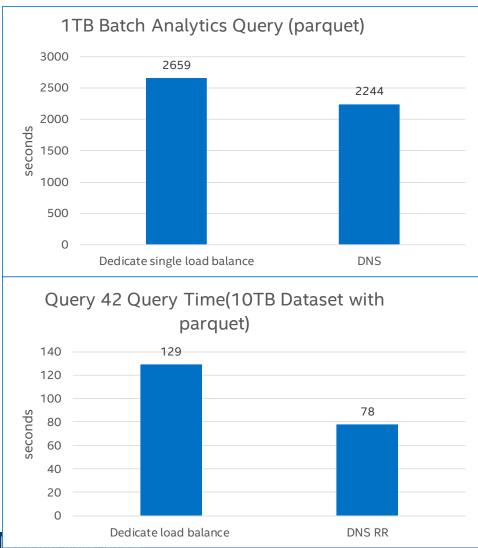
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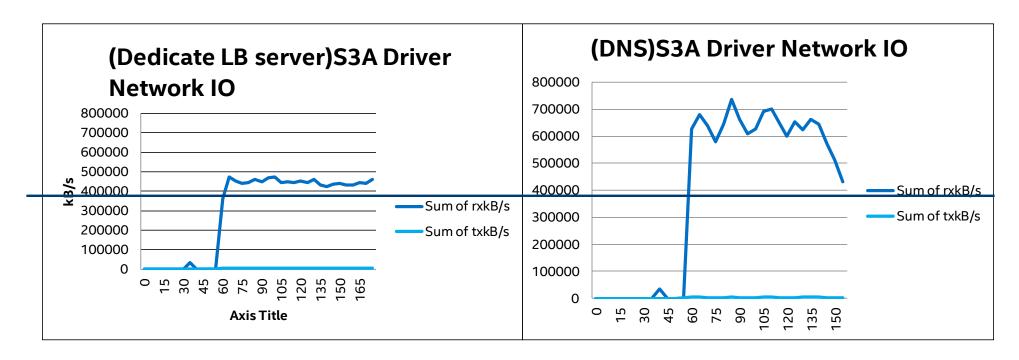
Performance evaluation --More RGWs and round-robin DNS



- 18% performance improvement with more RGWs and round-robin DNS
- Query42(has less shuffle) is 1.64x
 faster in the new architecture



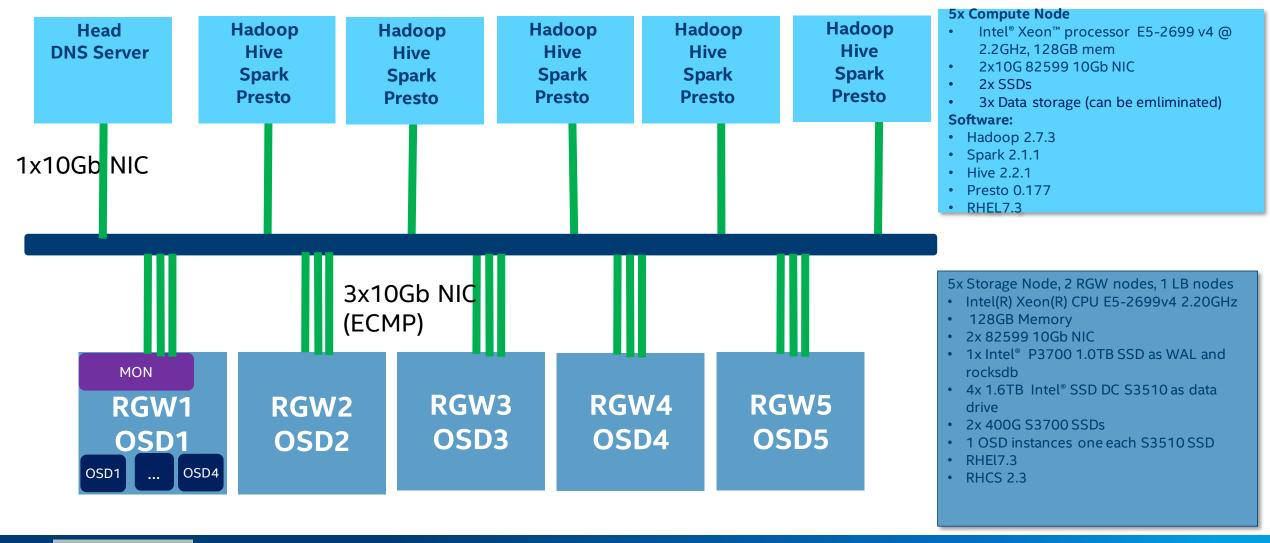
Key Resource Utilization Comparison



 Compute side(Hadoop s3a driver) can read more data from OSD faster, which showed DNS deployment bring big improvements for network throughput performance than single gateway with bonding/teaming technology

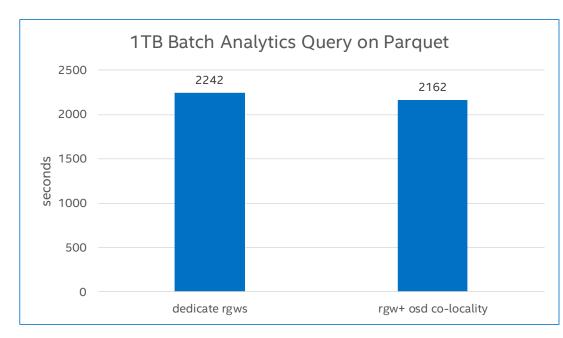


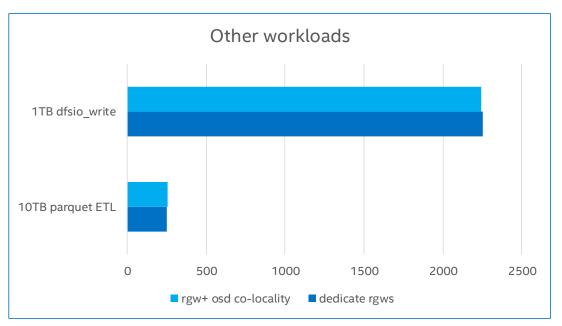
Hardware Configuration ---RGW and OSD Collocated





Performance under RGW & OSD Collocated

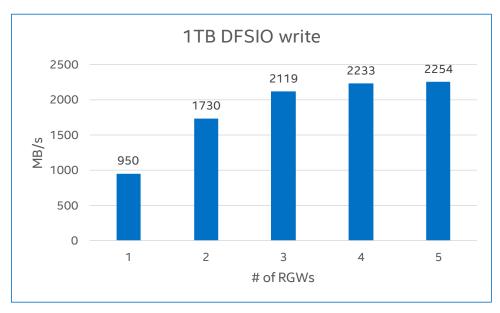


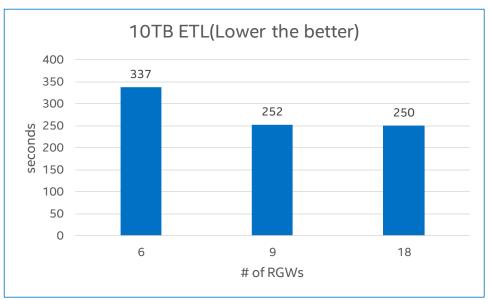


- No need extra dedicate RGW servers, RGW instance and OSD go through different network interface by enable ECMP
- No performance degradation, but less TCO



RGW & OSD collocated – RGW scaling



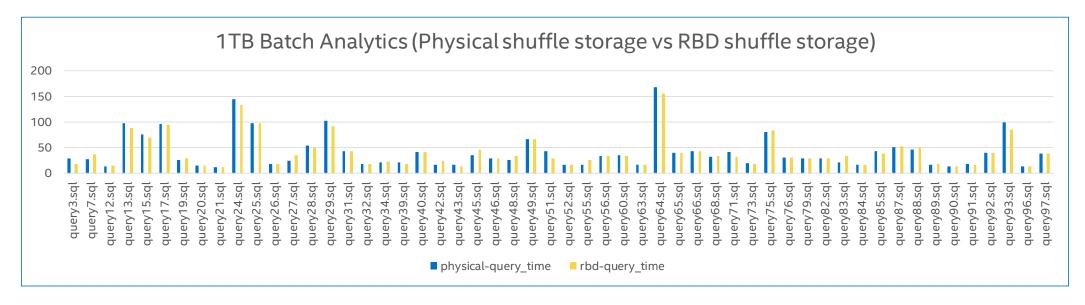


- Scale out RGWs can improve performance before OSD(storage) saturating
- So How many RGWs can win the best performance should be decided by the bandwidth of each RGW server and throughput of OSDs



Using Ceph RBD as shuffle Storage

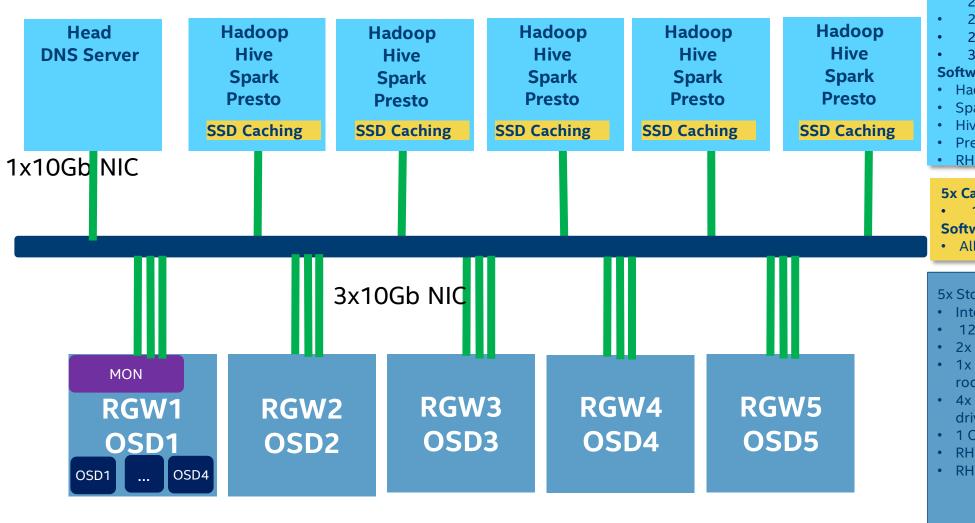
-- Eliminate the physical drive on the compute



- Remote RBD volumes on compute node to act as shuffle devices instead of physical shuffle device.
- For most queries the performance is not impacted.



Compute-side caching



5x Compute Node

- Intel® Xeon™ processor E5-2699 v4 @ 2.2GHz, 128GB mem
- 2x10G 82599 10Gb NIC
- 2x SSDs
- 3x Data storage (can be emliminated)

Software:

- Hadoop 2.7.3
- Spark 2.1.1
- Hive 2.2.1
- Presto 0.177
- RHEL7.3

5x Caching Node(co-located with compute)

1TB SSD(P3700)

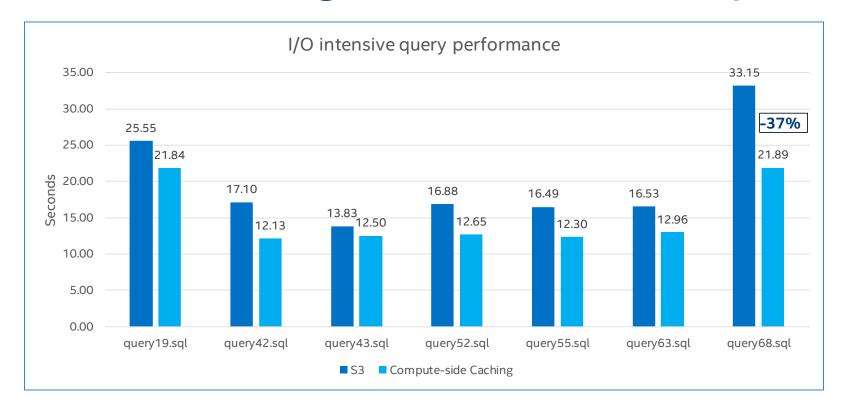
Software:

Alluxio* 1.7.0

5x Storage Node, 5 RGW nodes

- Intel(R) Xeon(R) CPU E5-2699v4 2.20GHz
- 128GB Memory
- 2x 82599 10Gb NIC
- 1x Intel® P3700 1.0TB SSD as WAL and rocksdb
- 4x 1.6TB Intel® SSD DC S3510 as data drive
- 1 OSD instances one each S3510 SSD
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Compute-side caching for I/O intensive queries



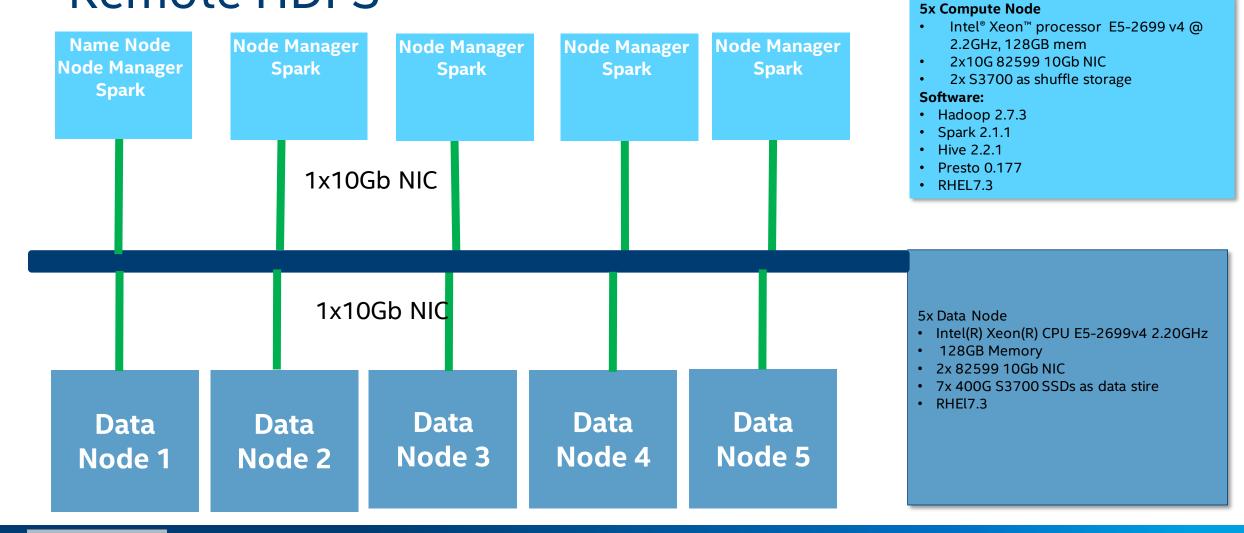
 Compute-side caching brings better efficiency(10% - 30%) for I/O intensive queries



PERFORMANCE COMPARISON WITH REMOTE HDFS



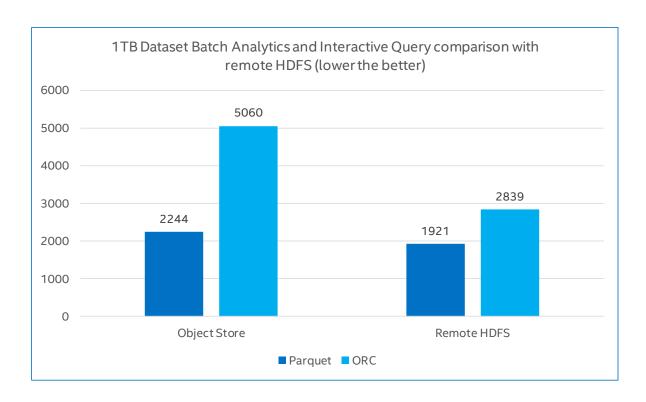
Hardware Configuration --Remote HDFS







Bigdata on Cloud vs. Remote HDFS --Batch Analytics



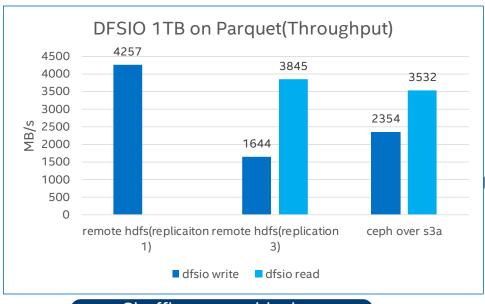
On-par performance compared with remote HDFS

- With optimizations, bigdata analytics on object storage is onpar with remote, especially on parquet format data
- performance of s3a driver close to native dfsclient, and demonstrate compute and storage separate solution has a considerable performance compare with combination solution



Bigdata on Cloud vs. Remote HDFS

--DFSIO



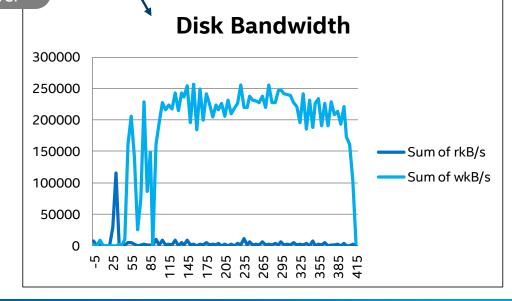
Device:	rrqm/s	wrqm/s	r/s	W/s	rkB/s	wkB/s	avgrq-sz a	vgqu-sz	await n	_await w	_await	svctm	%util
nvme0n1	0.00	0.00	0.00	4943.00	0.00	316696.00	128.14	0.27	0.06	0.00	0.06	0.03	13.25
sdb	0.00	0.00	5.00	100.50	20.00	15076.00	286.18	0.09	0.82	0.10	0.85	0.48	5.10
sdc	0.00	0.00	5.50	118.50	52.00	19969.50	322.93	0.13	1.03	0.09	1.08	0.62	7.70
sdh	0.00	0.00	4.50	120.50	108.00	17768.00	286.02	0.12	0.92	0.33	0.95	0.54	6.70
sda	0.00	0.00	23.50	474.00	274.00	66450.00	268.24	0.82	1.65	1.60	1.66	0.83	41.20
sdd	0.00	0.00	22.00	455.50	208.00	61194.00	257.18	2.59	5.43	4.89	5.45	1.94	92.70
sdg sdf	0.00	0.00	10.00	461.50	100.00	102906.00	436.93	83.39	176.86	15.95	180.35	2.04	96.15
sdf	0.00	43.50	19.50	922.50	168.00	110178.00	234.28	7.49	7.95	2.26	8.07	0.54	50.55
sdi	0.00	0.00	0.00	18.00	0.00	931.75	103.53	0.10	5.56	0.00	5.56	4.86	8.75
sde	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00

DFSIO write performance in ceph is better than remote hdfs(43%), but read performance is 34% lower

Write to Ceph hit disk bottleneck

Shuffle storage block at consuming data from Data Lake

ESTAB	2135680 0	10.0.2.16:8080	10.0.2.32:48046	timer:(keepalive,119min,0)
ESTAB	2135680 0	10.0.2.16:8080	10.0.2.32:47974	timer:(keepalive,119min,0)
ESTAB	2135680 0	10.0.2.16:8080	10.0.2.32:48006	timer:(keepalive,119min,0)
ESTAB	562816 0	10.0.2.16:8080	10.0.2.32:47956	timer:(keepalive,119min,0)
ESTAB	0 0	10.0.2.16:8080	10.0.2.32:47972	timer:(keepalive,119min,0)
ESTAB	2135680 0	10.0.2.16:8080	10.0.2.32:47978	timer:(keepalive,119min,0)
ESTAB	1611392 0	10.0.2.16:8080	10.0.2.32:48012	timer:(keepalive,119min,0)
ESTAB	2135680 0	10.0.2.16:8080	10.0.2.32:47942	timer:(keepalive,119min,0)
ESTAB	0 0	10.0.2.16:8080	10.0.2.32:47992	timer:(keepalive,119min,0)
ESTAB	0 0	10.0.2.16:8080	10.0.2.32:48030	timer:(keepalive,119min,0)
ESTAB	562816 0	10.0.2.16:8080	10.0.2.32:48004	timer:(keepalive,119min.0)



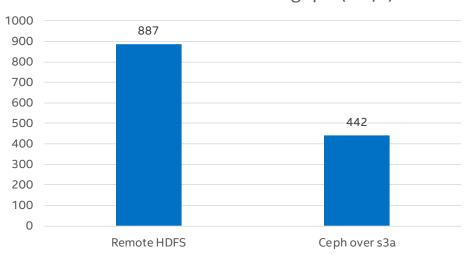


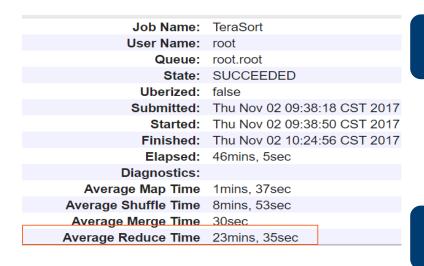


Bigdata on Cloud vs Remote HDFS

--Terasort

1TB Terasort Total Throughput(MB/s)

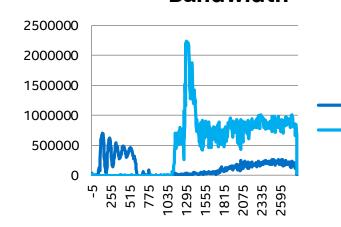




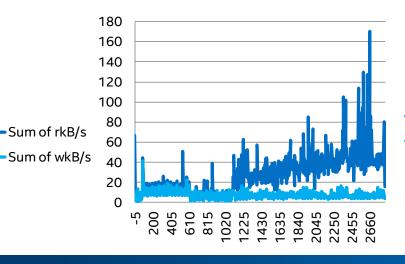
Time cost at Reduce stage is big part

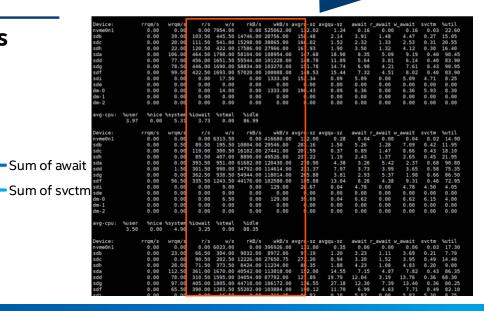
Read and write concurrently

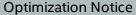
OSD Data Drive Disk Bandwidth



OSD Data Drive IO Latencies









Bigdata on Cloud vs. Remote HDFS --Ongoing rename optimizations

DirectOutputCommitter	An implementation in Spark 1.6, that return the destination address as working directory then no need to rename/move task output, no good robustness for failures, removed in Spark 2.0
IBM's "Stocator" committer	Targets Openstack Swift, good robustness, but it is another file system for s3a
Staging committer	A choice of new s3a committer, need large capacity of hard disk for staged data
Magic committer	A choice of new s3a committer, if you know your object store is consistent or use s3gurad, this committer has higher performance

"Renaming" overhead can be improved!



SUMMARY & NEXT STPES



Summary and Next Steps

Summary

- Bigdata on Ceph data lake is functionality ready validated by industry standard decision making workloads TPC-DS
- Bigdata on the Cloud delivers on-par performance with remote HDFS for batch analytics, intensive write operations still need further optimizations
- All flash solutions demonstrated significant TCO benefit compared with HDD solutions

Next

- Expand analytic workloads scope
- Rename operations optimizations to improve the performance
- Accelerating the performance with
 - speed up layer for shuffle
 - Compute-side caching







BACKUP



Experiment environment

Cluster	Hadoop head Hadoop slave		Load balancer	OSD	RGW	
Roles	Hadoop name node Secondary name node Resource manager Data node Node manager Hive metastore service Yarn history server Spark history server	Data node Node manager Presto server	Haproxy	Ceph osd	Ceph rados gateway	
# of node	1	5	1	5	5	
Processor	Intel(R)	Intel(R) Xeon(R) CPU E31280 @ 3.50GH 4 cores HT enabled				
Memory	128	3GB	128GB	128GB	32GB	
Storage		3 HDD SSD(vs s3700 metrics)	1x Intel S3510 480 GB SSD	 1x Intel® P3700 1.6TB as jounal4x 1.6TB Intel®SSD DC S3510 2X400GB s370 as data store 	1x Intel S3510 480 GB SSD	
Network	10	GB	40GB	10GB+10GB	10GB	



SW Configuration				
Hadoop version	2.7.3/2.8.1			
Spark version	2.1.1/2.2.0			
Hive version	2.2.1			
Presto version	0.177			
Executor memory	22GB			
Executor cores	5			
# of executor	24			
JDK version	1.8.0_131			
Memory.overhead	5GB			

S3A Key Performance Configuration				
fs.s3a.connection.maximu m	10			
fs.s3a.threads.max	30			
fs.s3a.socket.send.buffer	8192			
fs.s3a.socket.recv.buffer	8192			
fs.s3a.threads.keepalivetim e	60			
fs.s3a.max.total.tasks	1000			
fs.s3a.multipart.size	100M			
fs.s3a.block.size	32M			
fs.s3a.readahead.range	64k			
fs.s3a.fast.upload	true			
fs.s3a.fast.upload.buffer	array			
fs.s3a.fast.upload.active.bl ocks	4			
fs.s3a.experimental.input.f advise	radom			



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Computenode

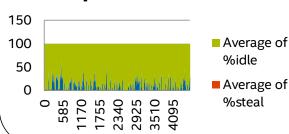
Resource Utilization on 1TB parquet

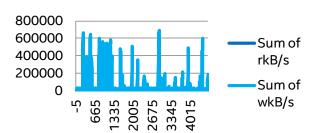
Cpu Utilization

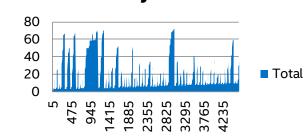
Disk Bandwidth

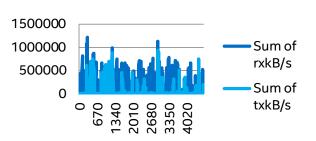
Memory Utilization



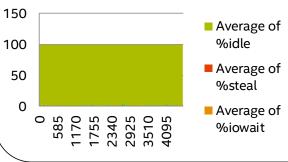




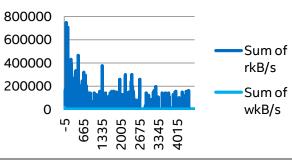




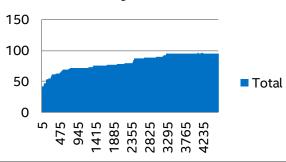
Cpu Utilization



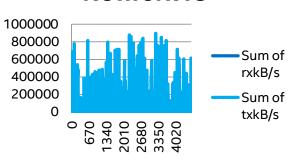




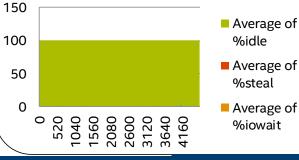
Memory Utilization



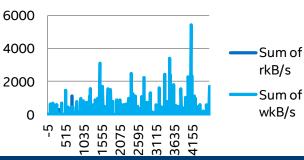
Network IO



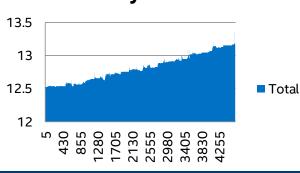
Cpu Utilization



Disk Bandwidth



Memory Utilization



Network IO

